

*Dual-guided pivot rules for LP*  
+  
*Vector Space Decomposition for LP*

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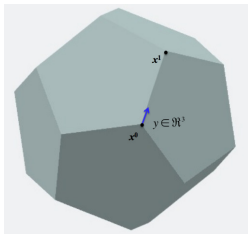
What are the links between PS, DCA, IPS, and MMCC ?  
... and Dantzig-Wolfe decomposition ?

## LP IN STANDARD FORM

$$\min \quad \mathbf{c}^T \mathbf{x}$$

$$\mathbf{Ax} = \mathbf{b} \quad [\pi]$$

$$\mathbf{x} \geq \mathbf{0}$$

FROM  $\mathbf{x}^0$  TO  $\mathbf{x}^1$ 

- ① Find a *potential* improving direction  $\mathbf{y}^0 \in \mathbb{R}^n$ .
- ② Determine step-size  $\rho^0 \in \mathbb{R}$ .
- ③ Compute  $\mathbf{x}^1 := \mathbf{x}^0 + \rho^0 \mathbf{y}^0$ .

$\bar{c}_j = 0, \forall j \in B$  : PRICING FOR  $j \in N$  (NON-BASIC VARIABLES)

Selection of an entering variable into basis  $\mathbf{A}_B$  relies on the minimum reduced cost of non-basic variables

$$\boldsymbol{\pi}^\top = \mathbf{c}_B^\top \mathbf{A}_B^{-1} \quad \bar{c}_j = c_j - \boldsymbol{\pi}^\top \mathbf{a}_j, \quad \forall j \in N.$$

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FIND THE MINIMUM REDUCED COST VALUE  $\mu$  (OPTIMAL IF  $\mu \geq 0$ )

max  $\mu$

$$\mu \leq c_j - \boldsymbol{\pi}^\top \mathbf{a}_j, \quad \forall j \in N \quad [y_j]$$

\*\*\*  $\mu$  is the smallest reduced cost (given  $\boldsymbol{\pi}$ ). \*\*\*

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EQUIVALENT TO FINDING A CONVEX COMBINATION OF NON-BASIC VARIABLES

$$\begin{aligned} \mu = \min \quad & \sum_{j \in N} \bar{c}_j y_j \\ & \sum_{j \in N} y_j = 1 \quad [\mu] \\ & y_j \geq 0, \quad \forall j \in N \end{aligned}$$

$\mathbf{y}^0$  IN THE PRIMAL SIMPLEX

Direction  $\mathbf{y}^0 \in \mathbb{R}^n$  :

the selected entering variable,  
the non-selected non-basic variables (they remain at 0),  
and the basic ones.

Step size computed such that  $\mathbf{x}^1 := \mathbf{x}^0 + \rho^0 \mathbf{y}^0 \geq \mathbf{0}$ .



LEFT MULTIPLICATION BY BASIS INVERSE  $A_B^{-1}$ 

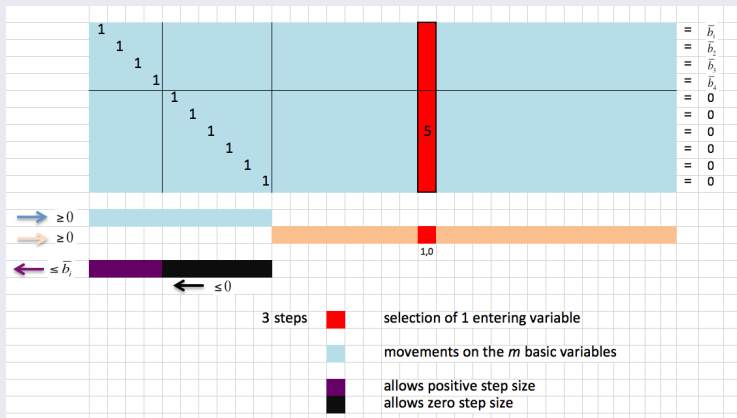
	$x_1$	$x_2$	$x_3$	$\lambda_4$	$\lambda_5$	$\lambda_6$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$		$\theta_w$
$c$	2	3	1				10	17	-20	14	-4		-5
	1						2	2	1	-5	7	=	30
		1					4	3	-5	10	-10	=	25
			1				-3	1	2	3	11	=	50
				1					6	5	-13	=	0
					1				3	4	-8	=	0
						1			3	-4	0	=	0
$x^0$	30	25	50	0	0	0	0	0	0	0	0	$185 = c^T x^0$	
							0.4	0.3	0.3				

$\psi$	
2	1
3	-2
1	5
?	0
?	0
?	0

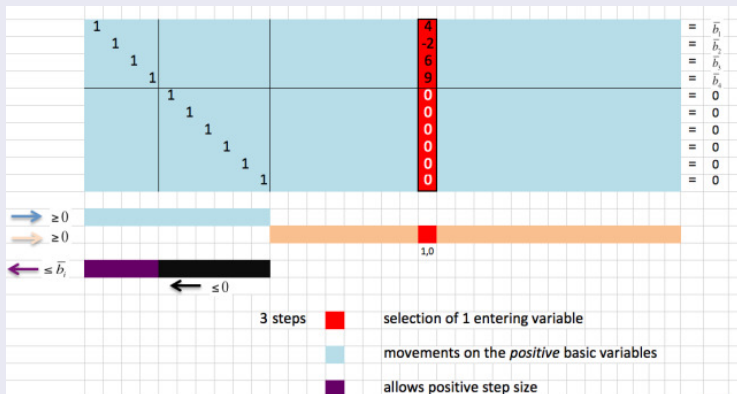


## PRIMAL SIMPLEX TABLEAU



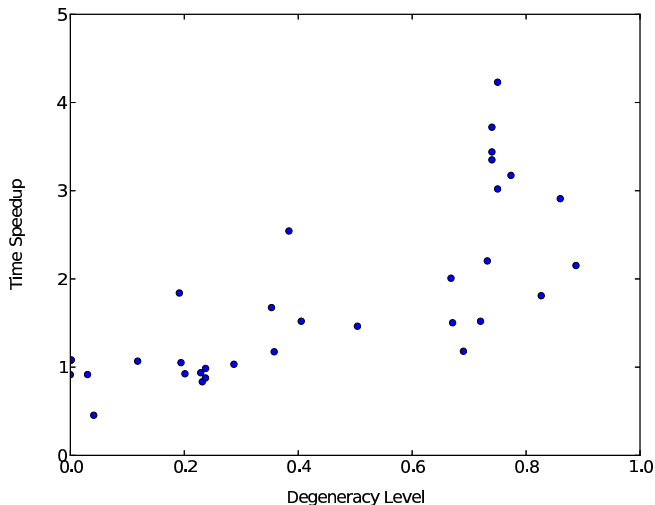
\*\*\* Changes on at most  $m + 1$  components. \*\*\*

## POSITIVE EDGE STRATEGY



\*\*\* Changes on at most  $p + 1$  components. \*\*\*  
Non-degenerate pivot.

SPEEDUP  $> 2$  FOR LPs WITH A DEGENERACY LEVEL ABOVE 25%



Vectors and matrices are written in **bold face**.

$\mathbf{I}_\ell$  : the  $\ell \times \ell$  identity matrix.

$\mathbf{0}$  ( $\mathbf{1}$ ) : a vector/matrix with all zeros (ones) entries of appropriate dimensions.

$\mathbf{A}_{RC}$  : sub-matrix of  $\mathbf{A}$  containing the rows and columns indexed by  $R$  and  $C$ .

Basis  $\mathbf{A}_B$ , inverse  $\mathbf{A}_B^{-1}$ ,  $\mathbf{c}_B^T \mathbf{x}_B$ ,  $\mathbf{A}_B \mathbf{x}_B$ ,  $\boldsymbol{\pi} = \mathbf{c}_B \mathbf{A}_B^{-1} \dots$

$\mathbf{l}_F < \mathbf{x}_F < \mathbf{u}_F$ ,  $\mathbf{x}_L = \mathbf{l}_L$ ,  $\mathbf{x}_U = \mathbf{u}_U$

USEFUL DECOMPOSITION OF  $\mathbf{x} \in \mathbb{R}^n$  IN  $\mathbf{Ax} = \mathbf{b}$ ,  $\mathbf{l} \leq \mathbf{x} \leq \mathbf{u}$

$$\mathbf{x} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix} \quad \mathbf{x} = \begin{bmatrix} \mathbf{x}_B \\ \mathbf{x}_N \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_{B_L} \\ \mathbf{x}_{B_U} \\ \mathbf{x}_{N_L} \\ \mathbf{x}_{N_U} \end{bmatrix}$$

$$\text{For } \emptyset \subseteq S \subseteq B : \mathbf{x} = \begin{bmatrix} \mathbf{x}_S \\ \mathbf{x}_{\bar{S}} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_{S_F} \\ \mathbf{x}_{S_L} \\ \mathbf{x}_{S_F} \\ \mathbf{x}_{\bar{S}_U} \\ \mathbf{x}_{\bar{S}_L} \\ \mathbf{x}_{\bar{S}_U} \end{bmatrix} \quad \mathbf{x} = \begin{bmatrix} \mathbf{x}_{\emptyset} \\ \mathbf{x} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix}$$

## GENERIC ALGORITHM WITH SINGLE PARAMETER SET $S$ , $\emptyset \subseteq S \subseteq B$

① Let  $k = 0$  and assume a feasible basic solution  $\mathbf{x}^k$  to LP (standard form).

② Select  $S \subseteq B$ . 
$$z^* := \min \quad \mathbf{c}_S^T \mathbf{x}_S + \mathbf{c}_{\bar{S}}^T \mathbf{x}_{\bar{S}}$$
  
 st. 
$$\mathbf{A}_S^T \mathbf{x}_S + \mathbf{A}_{\bar{S}}^T \mathbf{x}_{\bar{S}} = \mathbf{b}$$
  

$$\mathbf{x}_S, \mathbf{x}_{\bar{S}} \geq \mathbf{0}$$

③ From  $\mathbf{A}_S$ , retrieve an  $s \times s$  basis  $\mathbf{A}_{RS}$  for row-set  $R$ .

$$z^* := \min \quad \mathbf{c}_S^T \mathbf{x}_S + \mathbf{c}_{\bar{S}}^T \mathbf{x}_{\bar{S}}$$

$$\text{st.} \quad \begin{aligned} \mathbf{A}_{RS}^T \mathbf{x}_S + \mathbf{A}_{R\bar{S}}^T \mathbf{x}_{\bar{S}} &= \mathbf{b}_R & [\pi_R] \\ \mathbf{A}_{ZS}^T \mathbf{x}_S + \mathbf{A}_{Z\bar{S}}^T \mathbf{x}_{\bar{S}} &= \mathbf{b}_Z & [\pi_Z] \\ \mathbf{x}_S, \mathbf{x}_{\bar{S}} &\geq \mathbf{0} \end{aligned}$$

④ Fix  $\pi_R$  in row-set  $R$ ;  $\pi_Z$  is free.

⑤ Determine the smallest reduced cost  $\mu_S^k$ .

$$\mu_S^k := \max \mu \quad \text{st.} \quad \mu \leq c_j - \pi_R^T \mathbf{a}_{Rj} - \pi_Z^T \mathbf{a}_{Zj}, \quad \forall j$$

If  $\mu_S^k \geq 0$ , STOP. Current solution  $\mathbf{x}^k$  is optimal for LP.

⑥ Retrieve direction  $\mathbf{y}_S^k \in \mathbb{R}^n$  and compute its maximum step-size  $\rho_S^k$ .

⑦ Update  $\mathbf{x}^{k+1} := \mathbf{x}^k + \rho_S^k \mathbf{y}_S^k$ ;  $z^{k+1} := z^k + \rho_S^k \mu_S^k$ ;  $k := k + 1$ .

⑧ Goto Step 2.

LINEAR PROGRAM  $LP$

$$\begin{array}{ll} z^* := & \min \quad \mathbf{c}^T \mathbf{x} \\ \text{st.} & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{l} \leq \mathbf{x} \leq \mathbf{u} \end{array}$$

## LINEAR PROGRAM $LP$

$$\begin{aligned} z^* &:= \min \quad \mathbf{c}^T \mathbf{x} \\ \text{st.} \quad & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{l} \leq \mathbf{x} \leq \mathbf{u} \end{aligned}$$

## GENERIC ALGORITHM WITH SINGLE PARAMETER SET $S$ , $\emptyset \subseteq S \subseteq B$

- ① Let  $k = 0$  and assume a feasible basic solution  $\mathbf{x}^k$  to LP.
- ② Select  $\emptyset \subseteq S \subseteq B$ .
- ③ From  $\mathbf{A}_S$ , retrieve  $\mathbf{A}_{RS}$  and construct *residual problem*  $LP_S(\mathbf{x}^k)$ .
- ④ Fix dual variables in row-set  $R$ :  $\boldsymbol{\pi}_R$ .
- ⑤ Determine the value of the smallest reduced cost  $\mu_S^k$ .  
If  $\mu_S^k \geq 0$ , STOP. Current solution  $\mathbf{x}^k$  is optimal for LP.
- ⑥ Retrieve direction  $\mathbf{y}_S^k \in \mathbb{R}^n$  and compute its maximum step-size  $\rho_S^k$ .
- ⑦ Update  $\begin{aligned} \mathbf{x}^{k+1} &:= \mathbf{x}^k + \rho_S^k \mathbf{y}_S^k; \\ z^{k+1} &:= z^k + \rho_S^k \mu_S^k; \\ k &:= k + 1. \end{aligned}$
- ⑧ Goto Step 2.

ASSUME A FEASIBLE SOLUTION  $\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$

$$\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$$



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$$= \mathbf{x}^k + (\vec{\mathbf{y}} - \tilde{\mathbf{y}}),$$

$$\vec{\mathbf{y}}, \tilde{\mathbf{y}} \geq \mathbf{0}, \quad \vec{\mathbf{y}}^\top \tilde{\mathbf{y}} = 0, \quad \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k, \quad \tilde{\mathbf{y}} \leq \tilde{\mathbf{r}}^k$$

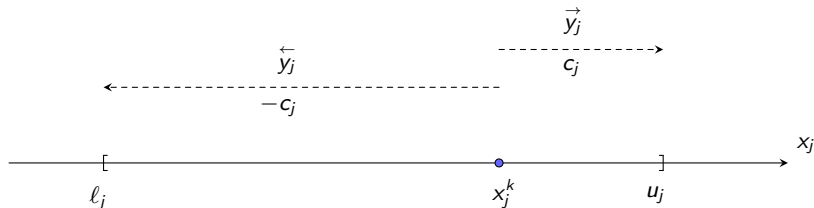
# STEP 3 CONSTRUCT RESIDUAL PROBLEM LP( $\mathbf{x}^k$ )

ASSUME A FEASIBLE SOLUTION  $\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$

$$\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$$

$$= \mathbf{x}^k + (\vec{\mathbf{y}} - \bar{\mathbf{y}}), \quad \vec{\mathbf{y}}, \bar{\mathbf{y}} \geq \mathbf{0}, \quad \vec{\mathbf{y}}^\top \bar{\mathbf{y}} = 0, \quad \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k, \quad \bar{\mathbf{y}} \leq \bar{\mathbf{r}}^k$$

$$= \mathbf{x}^k + \left( \begin{bmatrix} \vec{\mathbf{y}}_F \\ \vec{\mathbf{y}}_L \\ \vec{\mathbf{y}}_U \end{bmatrix} - \begin{bmatrix} \bar{\mathbf{y}}_F \\ \bar{\mathbf{y}}_L \\ \bar{\mathbf{y}}_U \end{bmatrix} \right); \quad \vec{\mathbf{y}}, \bar{\mathbf{y}} \geq \mathbf{0}, \quad \vec{\mathbf{y}}^\top \bar{\mathbf{y}} = 0, \quad \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k, \quad \bar{\mathbf{y}} \leq \bar{\mathbf{r}}^k$$



RESIDUAL PROBLEM  $LP(\mathbf{x}^k)$ ;  $\mathbf{x} := \mathbf{x}^k + (\vec{\mathbf{y}} - \tilde{\mathbf{y}})$  (CHANGE OF VARIABLES)

$$\begin{aligned}
z^* := \min \quad & \mathbf{c}^\top \mathbf{x}^k + \mathbf{c}^\top (\vec{\mathbf{y}} - \tilde{\mathbf{y}}) \\
\text{st.} \quad & \mathbf{A} \mathbf{x}^k + \mathbf{A}(\vec{\mathbf{y}} - \tilde{\mathbf{y}}) = \mathbf{b} \\
& \mathbf{0} \leq \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k \\
& \mathbf{0} \leq \tilde{\mathbf{y}} \leq \tilde{\mathbf{r}}^k
\end{aligned}$$

### STEP 3 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

RESIDUAL PROBLEM  $LP(\mathbf{x}^k)$ ;  $\mathbf{x} := \mathbf{x}^k + (\vec{\mathbf{y}} - \tilde{\mathbf{y}})$  (CHANGE OF VARIABLES)

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RESIDUAL PROBLEM  $LP(\mathbf{x}^k)$

$$\begin{aligned} z^* := \mathbf{c}^\top \mathbf{x}^k + \min \quad & \mathbf{c}^\top (\vec{\mathbf{y}} - \tilde{\mathbf{y}}) \\ \text{st.} \quad & \mathbf{A}(\vec{\mathbf{y}} - \tilde{\mathbf{y}}) = \mathbf{0} \\ & \mathbf{0} \leq \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k \\ & \mathbf{0} \leq \tilde{\mathbf{y}} \leq \tilde{\mathbf{r}}^k \end{aligned}$$

RESIDUAL PROBLEM WITH  $S = F$ ,  $\bar{S} = L \cup U$

$$z^* := \mathbf{c}^T \mathbf{x}^k +$$

$$\begin{array}{llllll} \min & \mathbf{c}_F^T(\vec{\mathbf{y}}_F - \tilde{\mathbf{y}}_F) & + & \mathbf{c}_L^T(\vec{\mathbf{y}}_L) & - & \mathbf{c}_U^T(\tilde{\mathbf{y}}_U) \\ \text{st.} & \mathbf{A}_F(\vec{\mathbf{y}}_F - \tilde{\mathbf{y}}_F) & + & \mathbf{A}_L(\vec{\mathbf{y}}_L) & - & \mathbf{A}_U(\tilde{\mathbf{y}}_U) = \mathbf{0} \\ & \vec{\mathbf{y}}_F \geq \mathbf{0}, \tilde{\mathbf{y}}_F \geq \mathbf{0}, & & \vec{\mathbf{y}}_L \geq \mathbf{0}, & & \tilde{\mathbf{y}}_U \geq \mathbf{0} \\ & \vec{\mathbf{y}}_F \leq \vec{\mathbf{r}}_F, \tilde{\mathbf{y}}_F \leq \tilde{\mathbf{r}}_F, & & \vec{\mathbf{y}}_L \leq \vec{\mathbf{r}}_L, & & \tilde{\mathbf{y}}_U \leq \tilde{\mathbf{r}}_U \end{array}$$

# STEP 3 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

RESIDUAL PROBLEM WITH  $S = B$  (BASIC),  $\bar{S} = N$  (NON-BASIC)

$$\begin{aligned}
 & \mathbf{c}^T \mathbf{x}^k + \min \mathbf{c}_B^T (\vec{\mathbf{y}}_B - \tilde{\mathbf{y}}_B) + \mathbf{c}_{N_L}^T \vec{\mathbf{y}}_{N_L} - \mathbf{c}_{N_U}^T \tilde{\mathbf{y}}_{N_U} \\
 \text{st. } & \mathbf{A}_B (\vec{\mathbf{y}}_B - \tilde{\mathbf{y}}_B) + \mathbf{A}_{N_L} \vec{\mathbf{y}}_{N_L} - \mathbf{A}_{N_U} \tilde{\mathbf{y}}_{N_U} = \mathbf{0} \\
 & \quad \quad \quad \vec{\mathbf{y}}_B \geq \mathbf{0}, \tilde{\mathbf{y}}_B \geq \mathbf{0} \quad \quad \quad \vec{\mathbf{y}}_{N_L} \geq \mathbf{0}, \tilde{\mathbf{y}}_{N_U} \geq \mathbf{0} \\
 & \quad \quad \quad \vec{\mathbf{y}}_F \leq \vec{\mathbf{r}}_F, \tilde{\mathbf{y}}_F \leq \tilde{\mathbf{r}}_F \\
 & \quad \quad \quad \vec{\mathbf{y}}_{B_L} \leq \mathbf{0}, \vec{\mathbf{y}}_{B_U} \leq \mathbf{0} \quad \quad \quad \vec{\mathbf{y}}_{N_L} \leq \vec{\mathbf{r}}_{N_L}, \tilde{\mathbf{y}}_{N_U} \leq \tilde{\mathbf{r}}_{N_U}
 \end{aligned}$$

$$B = F \cup B_L \cup B_U;$$

$$N = N_L \cup N_U$$

For  $\emptyset \subseteq S \subseteq B$ , find  $\mathbf{y}^k = \begin{bmatrix} \mathbf{y}_S^k \\ \mathbf{y}_{\bar{S}}^k \end{bmatrix} = \begin{bmatrix} (\bar{\mathbf{y}}_S^k - \bar{\mathbf{y}}_S^k) \\ (\bar{\mathbf{y}}_{\bar{S}}^k - \bar{\mathbf{y}}_{\bar{S}}^k) \end{bmatrix}$  of min reduced cost  $\mu^k$ .

For basic columns  $\mathbf{A}_S$ ,  
retrieve a partial basis  $\mathbf{A}_{RS}$ , a set of  $s$  independent rows.

$$\mathbf{T} = \begin{bmatrix} \mathbf{A}_{RS} & \mathbf{0} \\ \mathbf{A}_{ZS} & \mathbf{I}_{m-s} \end{bmatrix} \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_{RS}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZS}\mathbf{A}_{RS}^{-1} & \mathbf{I}_{m-s} \end{bmatrix}$$

For  $\emptyset \subseteq S \subseteq B$ , find  $\mathbf{y}^k = \begin{bmatrix} \mathbf{y}_S^k \\ \mathbf{y}_{\bar{S}}^k \end{bmatrix} = \begin{bmatrix} (\bar{\mathbf{y}}_S^k - \bar{\mathbf{t}}_S^k) \\ (\bar{\mathbf{y}}_{\bar{S}}^k - \bar{\mathbf{t}}_{\bar{S}}^k) \end{bmatrix}$  of min reduced cost  $\mu^k$ .

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Indeed,  $\mathbf{T}$  is a basis of  $\mathbb{R}^m$ .



## STEP 4 FIX DUAL VARIABLES IN ROW-SET $R$

$$\emptyset \subseteq S \subseteq B$$

$$\text{Perform } \mathbf{T}^{-1} \mathbf{A}(\vec{\mathbf{y}} - \tilde{\mathbf{y}}) = \mathbf{0}.$$

A block angular structure appears *after* the transformation by  $\mathbf{T}^{-1}$ .

GENERAL CASE

$$\pi^T = \psi^T \mathbf{T}^{-1}$$

$$\begin{aligned} \mathbf{c}^T \mathbf{x}^k + \min \quad & \mathbf{c}_S^T (\vec{\mathbf{y}}_S - \tilde{\mathbf{y}}_S) + \mathbf{c}_{\bar{S}}^T (\vec{\mathbf{y}}_{\bar{S}} - \tilde{\mathbf{y}}_{\bar{S}}) \\ \text{st.} \quad & (\vec{\mathbf{y}}_S - \tilde{\mathbf{y}}_S) + \bar{\mathbf{A}}_{R\bar{S}} (\vec{\mathbf{y}}_{\bar{S}} - \tilde{\mathbf{y}}_{\bar{S}}) = \mathbf{0} \quad [\psi_R = \mathbf{c}_S] \\ & \vec{\mathbf{y}}_S \geq \mathbf{0}, \tilde{\mathbf{y}}_S \geq \mathbf{0} \\ & \bar{\mathbf{A}}_{Z\bar{S}} (\vec{\mathbf{y}}_{\bar{S}} - \tilde{\mathbf{y}}_{\bar{S}}) = \mathbf{0} \quad [\psi_Z \text{ free}] \\ & \vec{\mathbf{y}}_{\bar{S}} \geq \mathbf{0}, \tilde{\mathbf{y}}_{\bar{S}} \geq \mathbf{0} \\ & \vec{\mathbf{y}}_S \leq \bar{\mathbf{r}}_S, \tilde{\mathbf{y}}_S \leq \bar{\mathbf{r}}_S \quad \vec{\mathbf{y}}_{\bar{S}} \leq \bar{\mathbf{r}}_{\bar{S}}, \tilde{\mathbf{y}}_{\bar{S}} \leq \bar{\mathbf{r}}_{\bar{S}} \end{aligned}$$

Observe  $\bar{\mathbf{c}}_S = \mathbf{0}$  (basic variables).

Master basis  $\mathbf{I}_S$ .      DW subproblem in red.

## STEP 5 FIND MINIMUM REDUCED COST $\mu_S^k$

### PRICING OF THE VARIABLES

$S \subseteq B$ , hence  $\bar{c}_S = 0$ .

Therefore pricing of  $\bar{y}_S$  and  $\bar{y}_S$  needed to get partial direction  $(\bar{y}_S^k - \bar{y}_S^k)$

followed by impact on  $(\bar{y}_S^k - \bar{y}_S^k)$  to complete direction  $y^k = \begin{bmatrix} y_S^k \\ y_S^k \end{bmatrix}$ .

### PRIMAL/DUAL FORMULATIONS OF THE PRICING

$$\psi_R^T = c_S^T A_{RS}^{-1}$$

$$\begin{aligned} \max \mu \quad \text{st.} \quad \mu \mathbf{1}^T &\leq c_S^T - \psi_R^T A_{RS} - \psi_Z^T \bar{A}_{ZS} & [\bar{y}_S] \\ \mu \mathbf{1}^T &\leq -(c_S^T - \psi_R^T A_{RS} - \psi_Z^T \bar{A}_{ZS}) & [\bar{y}_S] \end{aligned}$$

### A CONVEX COMBINATION OF THE VARIABLES $\bar{y}_S$ AND $\bar{y}_S$

$$\begin{aligned} \mu &= \min (c_S^T - \psi_R^T A_{RS})(\bar{y}_S - \bar{y}_S) \\ \text{st.} \quad \mathbf{1}^T \bar{y}_S + \mathbf{1}^T \bar{y}_S &= 1 & [\mu] \\ A_{ZS}(\bar{y}_S - \bar{y}_S) &= 0 & [\psi_Z] \\ \bar{y}_S, \bar{y}_S &\geq 0, \bar{y}_{S_L} \leq 0, \bar{y}_{S_U} \leq 0 \end{aligned}$$

\*\*\* Optimal solution :  $\mu_S^k$ ,  $\bar{y}_S^k$  and  $\bar{y}_S^k$ . \*\*\*

# STEP 6 RETRIEVE DIRECTION $\mathbf{y}_S^k$ OF MINIMUM REDUCED COST $\mu_S^k$

... COMPLETE DIRECTION  $\mathbf{y}_S^k$

Given  $\mu^k < 0$  and  $(\vec{\mathbf{y}}_S^k - \overleftarrow{\mathbf{y}}_S^k)$ , find impacts on  $(\vec{\mathbf{y}}_S^k - \overleftarrow{\mathbf{y}}_S^k)$ .

$$\begin{aligned} (\vec{\mathbf{y}}_S - \overleftarrow{\mathbf{y}}_S) + \bar{\mathbf{A}}_{R\bar{S}}(\vec{\mathbf{y}}_{\bar{S}}^k - \overleftarrow{\mathbf{y}}_{\bar{S}}^k) &= \mathbf{0} \\ \vec{\mathbf{y}}_S \geq \mathbf{0}, \overleftarrow{\mathbf{y}}_S \geq \mathbf{0} & \qquad \qquad \qquad \vec{y}_j \overleftarrow{y}_j = 0, \quad \forall j \in S_F \end{aligned}$$

DIRECTION  $\mathbf{y}_S^k$

$$\mathbf{y}_S^k = \begin{bmatrix} (\vec{\mathbf{y}}_S^k - \overleftarrow{\mathbf{y}}_S^k) \\ (\vec{\mathbf{y}}_{\bar{S}}^k - \overleftarrow{\mathbf{y}}_{\bar{S}}^k) \end{bmatrix} = \begin{bmatrix} -\bar{\mathbf{A}}_{R\bar{S}}(\vec{\mathbf{y}}_{\bar{S}}^k - \overleftarrow{\mathbf{y}}_{\bar{S}}^k) \\ (\vec{\mathbf{y}}_{\bar{S}}^k - \overleftarrow{\mathbf{y}}_{\bar{S}}^k) \end{bmatrix} : \begin{bmatrix} \text{Impact level (MP)} \\ \text{Pricing level (SP)} \end{bmatrix}$$

$\vec{\mathbf{y}}_S^k, \overleftarrow{\mathbf{y}}_S^k$  take positive and negative parts of  $-\bar{\mathbf{A}}_{R\bar{S}}(\vec{\mathbf{y}}_{\bar{S}}^k - \overleftarrow{\mathbf{y}}_{\bar{S}}^k)$ , respectively.

## STEP 6 COMPUTE MAXIMUM STEP-SIZE $\rho_S^k$

$$\rho \begin{bmatrix} \vec{y}_{S_F}^k \\ \vec{y}_{S_L}^k \\ \vec{y}_{S_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{S_F}^k \\ \vec{r}_{S_L}^k \\ \mathbf{0} \end{bmatrix} ; \quad \rho \begin{bmatrix} \vec{y}_{S_F}^k \\ \vec{y}_{S_L}^k \\ \vec{y}_{S_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{S_F}^k \\ \mathbf{0} \\ \vec{r}_{S_U}^k \end{bmatrix}$$

$$\rho \begin{bmatrix} \vec{y}_{S_F}^k \\ \vec{y}_{S_L}^k \\ \mathbf{0} \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{S_F}^k \\ \vec{r}_{S_L}^k \\ \mathbf{0} \end{bmatrix} ; \quad \rho \begin{bmatrix} \vec{y}_{S_F}^k \\ \mathbf{0} \\ \vec{y}_{S_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{S_F}^k \\ \mathbf{0} \\ \vec{r}_{S_U}^k \end{bmatrix}$$

10 out of 12 types of residual upper bonds to verify.  $\emptyset \subseteq S \subseteq B$

## STEP 6 COMPUTE MAXIMUM STEP-SIZE $\rho_S^k$

$$\rho \begin{bmatrix} \vec{y}_F^k \\ \vec{y}_{B_L}^k \\ \vec{y}_{B_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_F^k \\ \vec{r}_{B_L}^k \\ \mathbf{0} \end{bmatrix} ; \quad \rho \begin{bmatrix} \vec{y}_F^k \\ \vec{y}_{B_L}^k \\ \vec{y}_{B_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_F^k \\ \mathbf{0} \\ \vec{r}_{B_U}^k \end{bmatrix}$$

$$\rho \begin{bmatrix} \vec{y}_{N_L}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{N_L}^k \end{bmatrix} ; \quad \rho \begin{bmatrix} \vec{y}_{N_U}^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_{N_U}^k \end{bmatrix}$$

8 types of residual upper bonds to verify for Primal Simplex.  $S = B$

STEP 6 COMPUTE MAXIMUM STEP-SIZE  $\rho_S^k$

$$\rho \begin{bmatrix} \bar{\mathbf{y}}_P^k \\ \bar{\mathbf{y}}_Z^k \end{bmatrix} \leq \begin{bmatrix} \bar{\mathbf{r}}_P^k \\ \mathbf{0} \end{bmatrix}$$

Only 2 types of residual upper bounds to verify for PS in standard form.  
 $S = B = P \cup Z$  (positive and zero variables)

STEP 6 COMPUTE MAXIMUM STEP-SIZE  $\rho_S^k$

$$\rho \begin{bmatrix} \vec{y}_F^k \\ \vec{y}_L^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_F^k \\ \vec{r}_L^k \end{bmatrix}; \quad \rho \begin{bmatrix} \vec{y}_F^k \\ \vec{y}_U^k \end{bmatrix} \leq \begin{bmatrix} \vec{r}_F^k \\ \vec{r}_U^k \end{bmatrix}$$

4 types of strictly positive residual upper bonds to verify in MMCC.  $S = \emptyset$

STEP 6 COMPUTE MAXIMUM STEP-SIZE  $\rho_S^k$

$$\rho \left[ \vec{y}_F^k \right] \leq \left[ \vec{r}_F^k \right]; \quad \rho \left[ \vec{y}_F^k \right] \leq \left[ \vec{r}_F^k \right]$$

$$\rho \left[ \vec{y}_L^k \right] \leq \left[ \vec{r}_L^k \right]; \quad \rho \left[ \vec{y}_U^k \right] \leq \left[ \vec{r}_U^k \right]$$

4 types of strictly positive residual upper bonds to verify in IPS.  $S = F$



$$\mathbf{T} = [\mathbf{A}_B \quad \emptyset], \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_B^{-1} \\ \emptyset \end{bmatrix}$$

$$B = F \cup B_L \cup B_U; \quad N = N_L \cup N_U$$

$$\mathbf{T} = [\mathbf{A}_B \quad \emptyset], \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_B^{-1} \\ \emptyset \end{bmatrix}$$

$$B = F \cup B_L \cup B_U; \quad N = N_L \cup N_U$$

## PRIMAL SIMPLEX METHOD (DANTZIG 1945)

$$\begin{aligned} \mathbf{c}^\top \mathbf{x}^k + \min \quad & \mathbf{c}_B^\top (\vec{\mathbf{y}}_B - \tilde{\mathbf{y}}_B) \quad + \quad \mathbf{c}_{N_L}^\top \vec{\mathbf{y}}_{N_L} - \mathbf{c}_{N_U}^\top \tilde{\mathbf{y}}_{N_U} \\ \text{st.} \quad & (\vec{\mathbf{y}}_B - \tilde{\mathbf{y}}_B) \quad + \quad \bar{\mathbf{A}}_{N_L} \vec{\mathbf{y}}_{N_L} - \bar{\mathbf{A}}_{N_U} \tilde{\mathbf{y}}_{N_U} \quad = \quad \mathbf{0} \quad [\psi^\top = \mathbf{c}_B] \end{aligned}$$

$$\vec{\mathbf{y}}_B \geq \mathbf{0}, \quad \tilde{\mathbf{y}}_B \geq \mathbf{0}$$

$$\vec{\mathbf{y}}_{N_L} \geq \mathbf{0}, \quad \tilde{\mathbf{y}}_{N_U} \geq \mathbf{0}$$

$$\vec{\mathbf{y}}_F \leq \vec{\mathbf{r}}_F, \quad \tilde{\mathbf{y}}_F \leq \vec{\mathbf{r}}_F$$

$$\tilde{\mathbf{y}}_{B_L} \leq \mathbf{0}, \quad \tilde{\mathbf{y}}_{B_U} \leq \mathbf{0}$$

$$\vec{\mathbf{y}}_{N_L} \leq \vec{\mathbf{r}}_{N_L}, \quad \tilde{\mathbf{y}}_{N_U} \leq \vec{\mathbf{r}}_{N_U}$$

$$*\rho_B \geq 0*$$

## PROPERTIES OF PS

No equality constraints in the pricing problem.

Pricing contains • convex combination of the non-basic variables

• non-negativity restrictions (a cone).

Due to the step-size  $*\rho_B \geq 0*$ , possible degenerate pivots.

Oscillation of  $\mu_B$ ; it may even not converge towards 0.

$$\mathbf{T} = \begin{bmatrix} \mathbf{A}_{RF} & \mathbf{0} \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} \end{bmatrix}, \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_{RF}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZF}\mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} \end{bmatrix}.$$

## SPECIAL CASE #2 : $S = F \Rightarrow \bar{S} = L \cup U$

$$\mathbf{T} = \begin{bmatrix} \mathbf{A}_{RF} & \mathbf{0} \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} \end{bmatrix}, \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_{RF}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZF}\mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} \end{bmatrix}.$$

### IMPROVED PRIMAL SIMPLEX METHOD (ELHALLAOU ET AL. 2011)

$$\begin{aligned} & \mathbf{c}^T \mathbf{x}^k + \min \quad \mathbf{c}_F^T (\bar{\mathbf{y}}_F - \tilde{\mathbf{y}}_F) + \mathbf{c}_L^T \bar{\mathbf{y}}_L - \mathbf{c}_U^T \tilde{\mathbf{y}}_U \\ \text{st.} \quad & (\bar{\mathbf{y}}_F - \tilde{\mathbf{y}}_F) + \bar{\mathbf{A}}_{RL} \bar{\mathbf{y}}_L - \bar{\mathbf{A}}_{RU} \tilde{\mathbf{y}}_U = \mathbf{0} \quad [\psi_R = \mathbf{c}_F] \\ & \bar{\mathbf{y}}_F \geq \mathbf{0}, \tilde{\mathbf{y}}_F \geq \mathbf{0} \\ & \bar{\mathbf{A}}_{ZL} \bar{\mathbf{y}}_L - \bar{\mathbf{A}}_{ZU} \tilde{\mathbf{y}}_U = \mathbf{0} \quad [\psi_Z] \\ & \bar{\mathbf{y}}_L \geq \mathbf{0}, \tilde{\mathbf{y}}_U \geq \mathbf{0} \\ & \bar{\mathbf{y}}_F \leq \bar{\mathbf{r}}_F, \tilde{\mathbf{y}}_F \leq \bar{\mathbf{r}}_F, \quad \bar{\mathbf{y}}_L \leq \bar{\mathbf{r}}_L, \tilde{\mathbf{y}}_U \leq \bar{\mathbf{r}}_U \quad * \rho > 0 * \end{aligned}$$

### PROPERTIES OF IPS

$f$  equality constraints in the master problem,  $m - f$  in the pricing problem.

Non-degenerate pivots only ( $\rho_F > 0$ ).

$z^0 > z^1 > z^2 > \dots = z^*$  cost **strictly decreasing** at each iteration ( $\rho_F > 0$ ).

Oscillation of  $\mu_F$  but converging towards 0.

Select  $S = \emptyset$ .  $\mathbf{T} = \begin{bmatrix} \emptyset & \mathbf{I}_m \end{bmatrix}$ ,  $\mathbf{T}^{-1} = \begin{bmatrix} \emptyset \\ \mathbf{I}_m \end{bmatrix}$

Select  $S = \emptyset$ .  $\mathbf{T} = \begin{bmatrix} \emptyset & \mathbf{I}_m \end{bmatrix}$ ,  $\mathbf{T}^{-1} = \begin{bmatrix} \emptyset \\ \mathbf{I}_m \end{bmatrix}$

## MINIMUM MEAN CYCLE-CANCELING ALGORITHM ADAPTED FOR LP

$$\begin{aligned} & \mathbf{c}^T \mathbf{x}^k + \min \quad \mathbf{c}^T (\vec{\mathbf{y}} - \tilde{\mathbf{y}}) \\ & \text{st.} \quad \mathbf{A}(\vec{\mathbf{y}} - \tilde{\mathbf{y}}) = \mathbf{0} \quad \text{Directions} \\ & \quad \quad \vec{\mathbf{y}}, \tilde{\mathbf{y}} \geq \mathbf{0} \quad \text{in the cone} \\ & \quad \quad \vec{\mathbf{y}}_L, \vec{\mathbf{y}}_U \leq \mathbf{0} \quad \text{at vertex } \mathbf{x}^k \\ & \quad \quad \vec{\mathbf{y}}_F, \tilde{\mathbf{y}}_F, \vec{\mathbf{y}}_L, \tilde{\mathbf{y}}_U \leq \vec{\mathbf{r}}_F^k, \tilde{\mathbf{r}}_F^k, \vec{\mathbf{r}}_L^k, \tilde{\mathbf{r}}_U^k \quad * \text{Step size } \rho_0 > 0 * \end{aligned}$$

## PROPERTIES OF MMCC

All equality constraints in the pricing problem.

Upper bounds in the master problem.

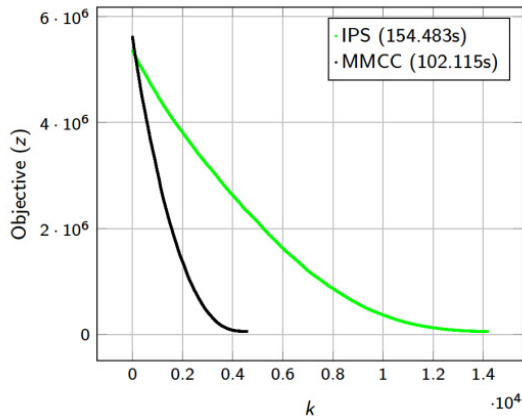
$z^0 > z^1 > z^2 > \dots = z^*$  cost **strictly decreasing** at each iteration ( $\rho_0 > 0$ ).

$\mu^0 \leq \mu^1 \leq \mu^2 \leq \dots = 0$  smallest reduced cost **non decreasing**.

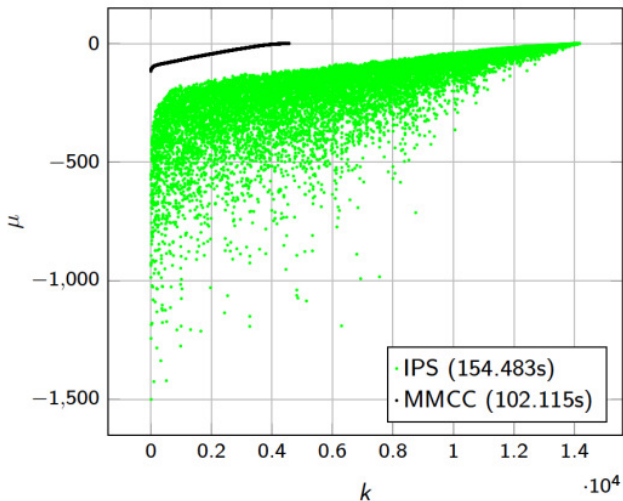
MMCC is strongly polynomial for network flow problems in  $O(mn)$  phases.

Goldberg and Tarjan (1989), Radzick and Goldberg (1994)

## COMPUTATIONAL RESULTS

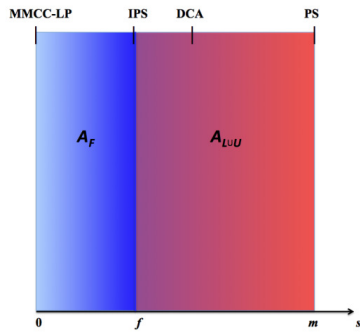


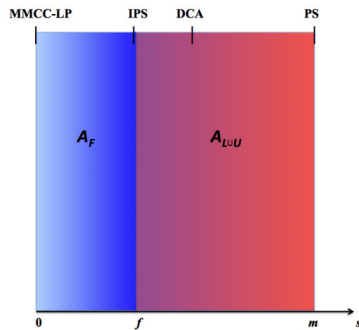
## COMPUTATIONAL RESULTS





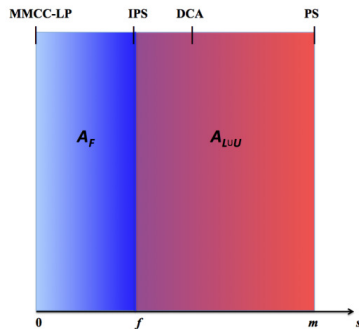
# SOME PROPERTIES



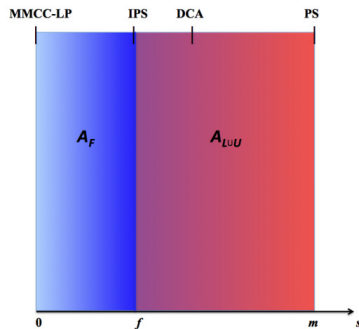


- $S \cap \{L \cup U\} \neq \emptyset$  : It may come up with degenerate pivots and not converge.
- Primal Simplex method (PS).

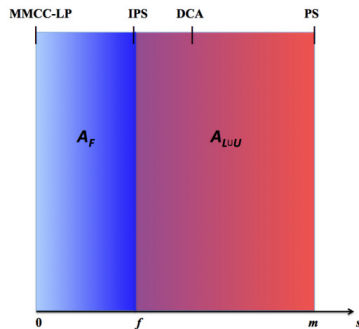
$$S = B, * \rho_B \geq 0 *$$



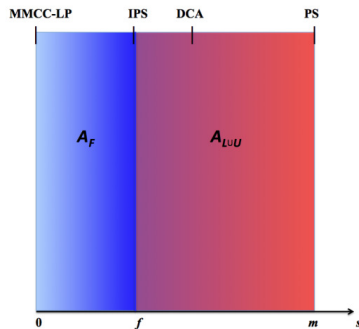
- $S \cap \{L \cup U\} \neq \emptyset$  : It may come up with degenerate pivots and not converge.
  - Primal Simplex method (PS).  $S = B, * \rho_B \geq 0 *$
- $\emptyset \subseteq S \subseteq F$  : It ensures a non-degenerate pivot at every iteration.
  - Improved Primal Simplex algorithm (IPS).  $S = F, * \rho_F > 0 *$



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  - Minimum mean cycle-canceling algorithm (MMCC)  $S = \emptyset, * \rho_\emptyset > 0 *$



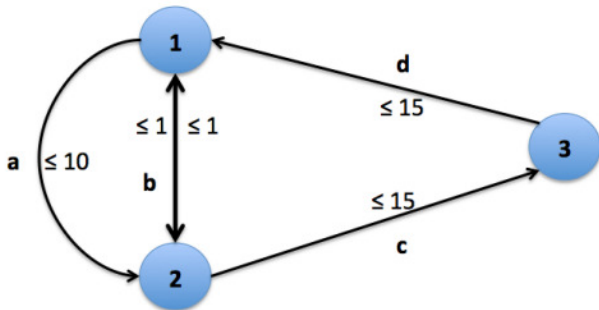
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*Strongly polynomial for network flow problems.*



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  - Minimum mean cycle-canceling algorithm (MMCC)  $S = \emptyset, * \rho_\emptyset > 0 *$   
*Strongly polynomial for network flow problems.*
  - \*  $S \subset F$  : Optimal direction  $y_s^k$  can be an interior ray.

# MMCC:

*non-extreme ray* a-c-d



\*\*\* Imagine the same transformation is kept for a while...

- $\mathbf{T} := [\mathbf{\Lambda}, \mathbf{\Lambda}^\perp]$ , where  $\mathbf{\Lambda}$  is a set of independent columns.

- $\mathbf{\Lambda} = \begin{bmatrix} \mathbf{\Lambda}_R \\ \mathbf{\Lambda}_Z \end{bmatrix}$ , where  $\mathbf{\Lambda}_R$  is a set of  $s$  independent rows of  $\mathbf{\Lambda}$ .

- $\mathbf{T} = \begin{bmatrix} \mathbf{\Lambda}_R & \mathbf{0} \\ \mathbf{\Lambda}_Z & \mathbf{I}_{m-s} \end{bmatrix} \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{\Lambda}_R^{-1} & \mathbf{0} \\ \mathbf{\Lambda}_Z \mathbf{\Lambda}_R^{-1} & \mathbf{I}_{m-s} \end{bmatrix}$

- $\mathbf{T}^{-1}$  splits row-space  $\mathbb{R}^m$  of  $LP(\mathbf{x}^k)$  into two vector subspaces  $\mathbf{V}$  and  $\mathbf{V}^\perp$ .

- Vector subspace basis  $\mathbf{\Lambda}$  spans  $\mathbf{V}$  of dimension  $0 \leq s \leq m$ .

- Vector  $\mathbf{a} \in \mathbf{V}$  if and only if  $\bar{\mathbf{a}}_Z = \mathbf{0}$ , where  $\bar{\mathbf{a}} = \mathbf{T}^{-1}\mathbf{a} = \begin{bmatrix} \bar{\mathbf{a}}_R \\ \mathbf{0} \end{bmatrix}$ .

- $S \subseteq B$  : index set of basic columns spanned by  $\mathbf{\Lambda}$ .

- Algorithmic properties derived according to subset  $S$ .



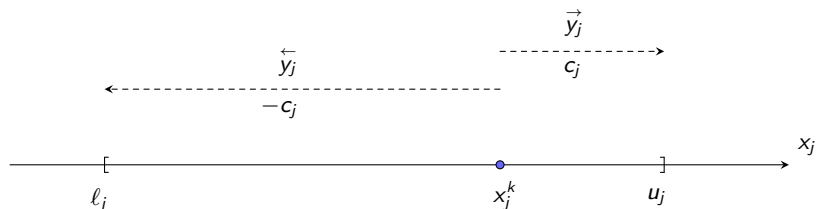
# DYNAMIC CONSTRAINT AGGREGATION FOR SET PARTITIONING (ELHALLAOUI ET AL. 2005)

The partition of the row-set is derived from the  $f$  groups of identical rows of  $A_F$ .

$$\begin{array}{c} \mathbf{A}_F \\ \left( \begin{array}{cccc} 1 & 1 & & \\ & 1 & 1 & \\ & & 1 & 1 \\ 1 & & 1 & \\ 1 & & & 1 \\ 1 & 1 & & \\ & 1 & 1 & \\ 1 & & & 1 \end{array} \right) \end{array} \mapsto \begin{array}{c} \mathbf{\Lambda} \\ \left( \begin{array}{cccc} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \\ \hline 1 & & & 1 \\ & 1 & & \\ & & & 1 \end{array} \right) \end{array} .$$

## PROPERTIES

$F \subseteq S \subseteq B$  for fractional solutions but  $S = F$  (or  $S = \emptyset$ ) for binary solutions.  
 Degenerate pivots may occur,  $\rho \geq 0$ , if  $r > f$ .

RESIDUAL PROBLEM  $LP(\mathbf{x}^k)$ 

$$\begin{aligned}
 z^* := \mathbf{c}^\top \mathbf{x}^k + \quad & \min && \mathbf{c}^\top (\vec{\mathbf{y}} - \tilde{\mathbf{y}}) \\
 \text{st.} &&& \mathbf{A}(\vec{\mathbf{y}} - \tilde{\mathbf{y}}) = \mathbf{0} \\
 &&& \mathbf{0} \leq \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k \\
 &&& \mathbf{0} \leq \tilde{\mathbf{y}} \leq \tilde{\mathbf{r}}^k
 \end{aligned}$$